



PROGRAMMING AND OPTIMIZATION FOR INTEL[®] ARCHITECTURE

The Hands-On Workshop (HOW) Series
Session 2

Colfax International — colfaxresearch.com

September 2016

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- ▶ HOW to Program Intel Architecture
 - 01. Parallelism, specialization, guided tour – Sep 26
 - 02. Programming Intel Xeon Phi (KNC, KNL) – Sep 27
- ▶ HOW to Express Parallelism
 - 03. Automatic vectorization – Sep 28
 - 04. Multi-threading with OpenMP – Sep 29
- ▶ HOW to Get Performance
 - 05. Comprehensive demo – Sep 30
 - 06. Scalar & vectorization tuning – Oct 3
 - 07. Multi-threading: common issues – Oct 4
 - 08. Multi-threading: memory aspect – Oct 5
 - 09. Memory traffic – Oct 6
- ▶ HOW to Scale
 - 10. Distributed Computing: MPI – Oct 7

September 2016						
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■ — Webinar+remote access						

Course page: colfaxresearch.com/how-16-09

- ▶ Slides (including this one), code downloads
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- ▶ Chat (during webinars or offline)



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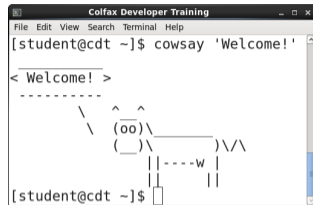
Welcome aboard. Post questions today!

Email (organizational):

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HANDS-ON EXERCISES AND REMOTE ACCESS

- ▶ 96 people receive a remote access token
- ▶ Virtualized Intel Xeon CPU, real Intel Xeon Phi coprocessor (1st gen, KNC), SW tools
- ▶ Can access the system the entire 2 weeks of the workshop



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Colfax Developer Training
File Edit View Search Terminal Help
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[student@cdt ~]$
```

- ▶ Not among the 96? Stay tuned: follow along with instructor, use own system, or wait for a seat
- ▶ Use it or lose it: if you do not log in for a while, remote access token goes to next student on the list



§2. PROGRAMMING COPROCESSORS



OFFLOAD AND NATIVE MODELS

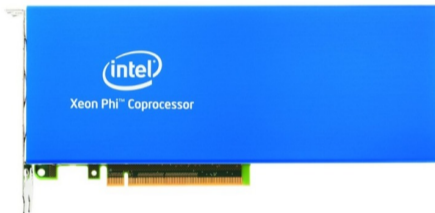
Intel Xeon Processor



Current: Broadwell
Upcoming: Skylake

Multi-Core Architecture

Intel Xeon Phi Coprocessor, 1st generation



Knights Corner (KNC)

Intel Xeon Phi Processor, 2nd generation*



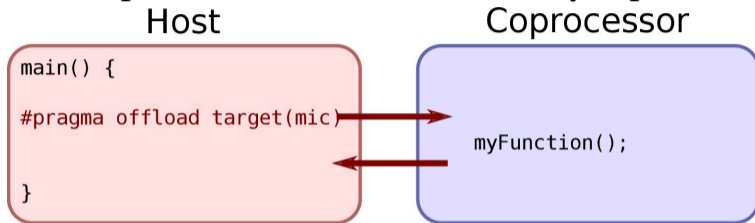
* socket and coprocessor versions

Knights Landing (KNL)

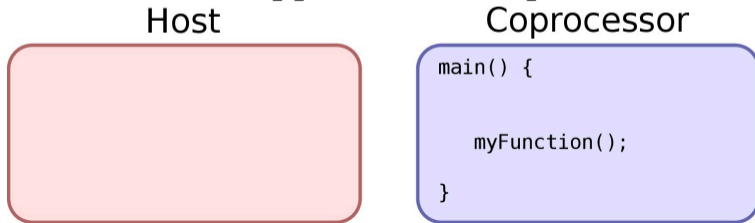
Intel Many Integrated Core (MIC) Architecture

OFFLOAD AND NATIVE MODELS

- ▶ Offload model (explicit/virtual-shared memory/OpenMP 4.0):



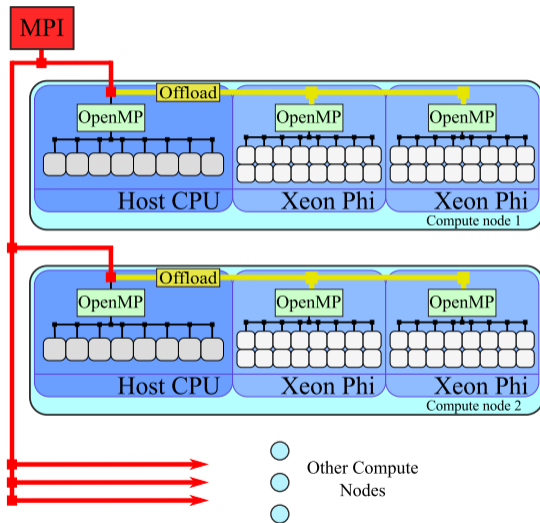
- ▶ Native model (standalone application/MPI process):



HETEROGENEOUS DISTRIBUTED COMPUTING WITH XEON PHI

Option 1: MPI+OpenMP with Offload.

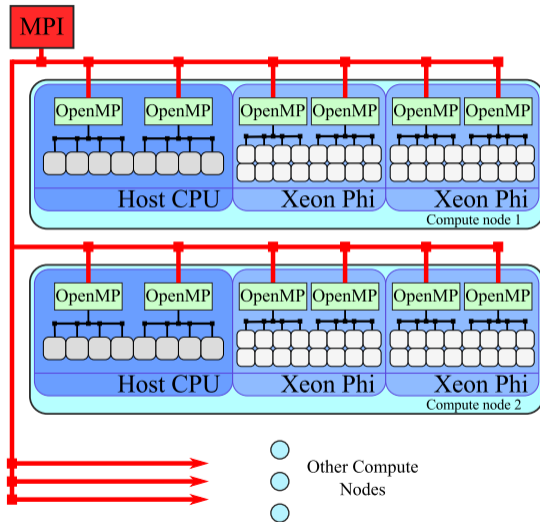
- ▶ MPI processes are multi-threaded with OpenMP.
- ▶ MPI runs only on CPUs.
- ▶ MPI processes offload to coprocessor(s).
- ▶ OpenMP in offload regions.



HETEROGENEOUS DISTRIBUTED COMPUTING WITH XEON PHI

Option 2: Symmetric hybrid MPI+OpenMP.

- ▶ MPI processes on hosts
- ▶ Native MPI processes on the coprocessor.
- ▶ Multi-threading with OpenMP.



OFFLOAD OR NATIVE: HOW TO DECIDE

Native	Offload
≤ 16 GiB	> 16 GiB
All parallel	Parallel + serial phases
Complex data structures	Bitwise-copyable data
Any arithmetic intensity	(FLOPs/transfer) \gg 1000

Native = same code on CPU and MIC

Offload = must insert directives in code



EXPLICIT OFFLOAD (LEO)

EXPLICIT OFFLOAD: PRAGMA-BASED APPROACH

“Hello World” in the explicit offload model:

```
1 #include <stdio>
2 int main() {
3     printf("Hello World from host!\n");
4     #pragma offload target(mic)
5     {
6         printf("Hello World from coprocessor!\n"); fflush(stdout);
7     }
8     printf("Bye\n");
9 }
```

Application runs on the host, but some parts of code and data are moved (“offloaded”) to the coprocessor.

Detailed syntax in the [Intel C++ Compiler Reference](#).

COMPILING AND RUNNING AN OFFLOAD APPLICATION

```
vega@lyra% icpc hello_offload.cc -o hello_offload
vega@lyra% ./hello_offload
Hello World from host!
Bye
Hello World from coprocessor!
```

- ▶ No additional arguments (for Intel compiler)
- ▶ Launch on host as a regular application
- ▶ Code inside of `#pragma offload` is offloaded automatically
- ▶ Console output on coprocessor buffered, mirrored to the host
- ▶ If no coprocessor available, default behavior is error; may be overridden to fall back to host



OFFLOADING FUNCTIONS AND DATA

OFFLOADING FUNCTIONS

```
1 __attribute__((target(mic))) void MyFunction() {  
2     // ... implement function as usual  
3 }  
4  
5 int main(int argc, char * argv[] ) {  
6     #pragma offload target(mic)  
7     {  
8         MyFunction();  
9     }  
10 }
```

- ▶ Functions used on coprocessor must be marked with the specifier `__attribute__((target(mic)))`
- ▶ Compiler produces a host version and a coprocessor version of such functions (to enable fall-back to host)

OFFLOADING MULTIPLE FUNCTIONS

```
1 #pragma offload_attribute(push, target(mic))
2 void MyFunctionOne() {
3 // ... implement function as usual
4 }
5
6 void MyFunctionTwo() {
7 // ... implement function as usual
8 }
9 #pragma offload_attribute(pop)
```

- ▶ To mark a long block of code with the offload attribute, use `#pragma offload_attribute(push/pop)`

OFFLOADING DATA: LOCAL SCALARS AND ARRAYS

```
1 void MyFunction() {  
2     const int N = 1000;  
3     int data[N];  
4     #pragma offload target(mic)  
5     {  
6         for (int i = 0; i < N; i++)  
7             data[i] = 0;  
8     }
```

- ▶ Scope-local scalars and known-size arrays offloaded automatically
- ▶ Data is copied from host to coprocessor at the start of offload
- ▶ Data is copied back from coprocessor to host at the end of offload
- ▶ Bitwise-copyable data only (arrays of basic types and scalars)
C++ classes, etc. should use virtual-shared memory model

OFFLOADING DATA: GLOBAL AND STATIC VARIABLES

```
1 int* __attribute__((target(mic))) data;
2
3 void MyFunction() {
4     static int __attribute__((target(mic))) N;
5     // ...
6 }
7
8 int main() {
9     // ...
10 }
```

- ▶ Global and static variables must be marked with the offload attribute
- ▶ `#pragma offload_attribute(push/pop)` may be used as well

DATA MARSHALLING FOR DYNAMICALLY ALLOCATED DATA

```
1 double *p1=(double*)malloc(sizeof(double)*N);  
2 double *p2=(double*)malloc(sizeof(double)*N);  
3  
4 #pragma offload target(mic) in(p1 : length(N)) out(p2 : length(N))  
5 {  
6     // ... perform operations on p1[] and p2[]  
7 }
```

- ▶ #pragma offload recognizes clauses in, out, inout and nocopy
- ▶ Data size (length), alignment, redirection, and other properties may be specified
- ▶ Marshalling is required for pointer-based data

OPTIONAL OFFLOAD, FALL-BACK TO HOST

```
1 #pragma offload target(mic) optional
2 {
3     printf("Hello World! I have %d logical processors.\n",
4         sysconf(_SC_NPROCESSORS_ONLN )); fflush(stdout);
5 }
```

```
vega@lyra% icpc Offload-Fallback.cc -o Offload-Fallback
vega@lyra% ./Offload-Fallback
Hello World! I have 244 logical processors.
vega@lyra% sudo systemctl stop mpss # Disabling coprocessors
vega@lyra% ./Offload-Fallback
Hello World! I have 48 logical processors.
```



MULTIPLE COPROCESSORS WITH EXPLICIT OFFLOAD

MULTIPLE COPROCESSORS WITH EXPLICIT OFFLOAD

▶ Querying the number of coprocessors:

```
1  const int numDevices = _Offload_number_of_devices();  
2  printf("Number of available coprocessors: %d\n" , numDevices);
```

▶ Specifying offload target:

```
1  #pragma offload target(mic: 0)  
2  { /* ... */ }
```

▶ Query the device number from within Offload:

```
1  #pragma offload target(mic)  
2  {  
3    const int deviceNum = _Offload_get_device_number();  
4    printf("Hello from coprocessor %d!\n" , deviceNum);  
5  }
```

MULTIPLE OFFLOADS USING HOST THREADS

```
1  const int nDevices = _Offload_number_of_devices();
2  #pragma omp parallel num_threads(nDevices)
3  {
4      const int i = omp_get_thread_num();
5      #pragma offload target(mic: i)
6          {
7              MyFunction(/*...*/);
8          }
9  }
```

- ▶ Up to 8 coprocessors, up to 56 host threads
- ▶ All offloads start simultaneously and block the respective thread

DYNAMIC WORK DISTRIBUTION

```
1  const int nDevices = _Offload_number_of_devices();
2  omp_set_num_threads(nDevices);
3  #pragma omp parallel for schedule(dynamic, 1)
4      for (int i = 0; i < nWorkItems; i++) {
5          const int iDevice = omp_get_thread_num();
6          #pragma offload target(mic: iDevice)
7              {
8                  MyFunction(i);
9              }
10 }
```

- ▶ Up to 8 coprocessors, up to 32 host threads
- ▶ nWorkItems are dynamically scheduled on nDevices



MEMORY ALLOCATION CONTROL

MEMORY RETENTION AND DATA PERSISTENCE ON COPROCESSOR

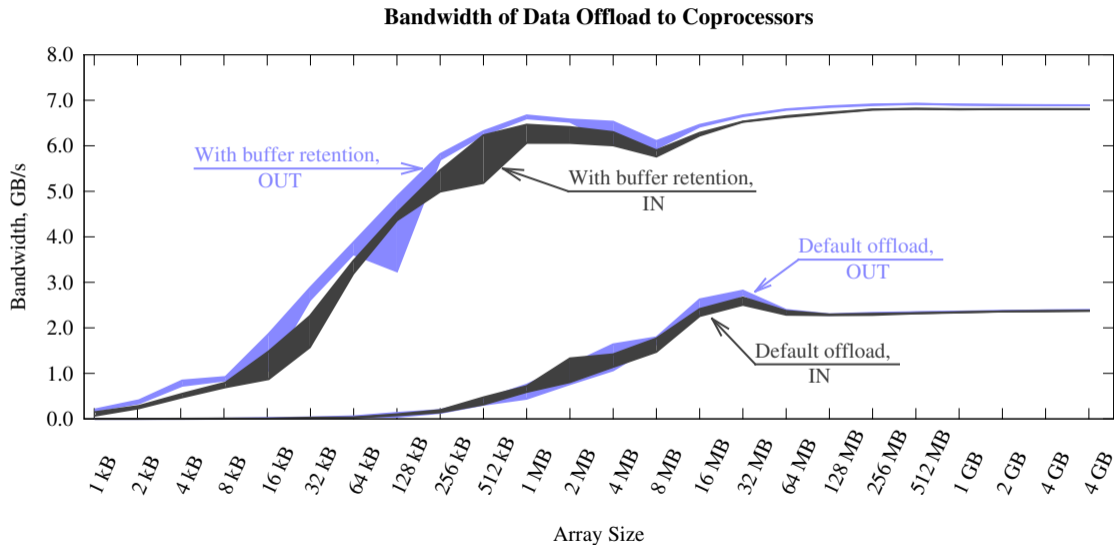
- ▶ By default, memory on coprocessor is allocated before, deallocated after offload
- ▶ Specifiers `alloc_if` and `free_if` allow to avoid allocation/deallocation
- ▶ Data transfer across the PCIe bus rate is ≈ 7 GB/s
- ▶ To allocate memory on the coprocessor – 0.5-2.0 GB/s

```

1 #pragma offload target(mic:0) in(p : length(N) alloc_if(1) free_if(0) )
2 { /* allocate memory for array p on coprocessor, do not deallocate */ }
3
4 #pragma offload target(mic:0) in(p : length(N) alloc_if(0) free_if(0) )
5 { /* re-use previously allocated memory buffer on coprocessor */ }
6
7 #pragma offload target(mic:0) in(p : length(0) alloc_if(0) free_if(0) )
8 { /* re-use previously transferred data on coprocessor */ }
9
10 #pragma offload target(mic:0) out(p : length(N) alloc_if(0) free_if(1) )
11 { /* re-use memory and deallocate at the end of offload */ }

```

OFFLOAD LATENCY WITH AND WITHOUT MEMORY/DATA RETENTION



PRECAUTIONS WITH PERSISTENT DATA

- ▶ Use explicit zero-based coprocessor number (e.g., `mic:0` as shown below)
- ▶ With multiple coprocessors, if target number is unspecified, any coprocessor can be used, which will result in runtime errors if persistent data cannot be found.

```
1 #pragma offload target(mic:0) in(p : length(N) alloc_if(1) free_if(0) )  
2 { /* allocate memory for array p on coprocessor, do not deallocate */ }
```

- ▶ Do not change the value of the host pointer to a persistent array: the runtime system finds the data on coprocessor using the host pointer value, not variable name.



OVERLAPPING COMMUNICATION AND COMPUTATION

ASYNCHRONOUS OFFLOAD

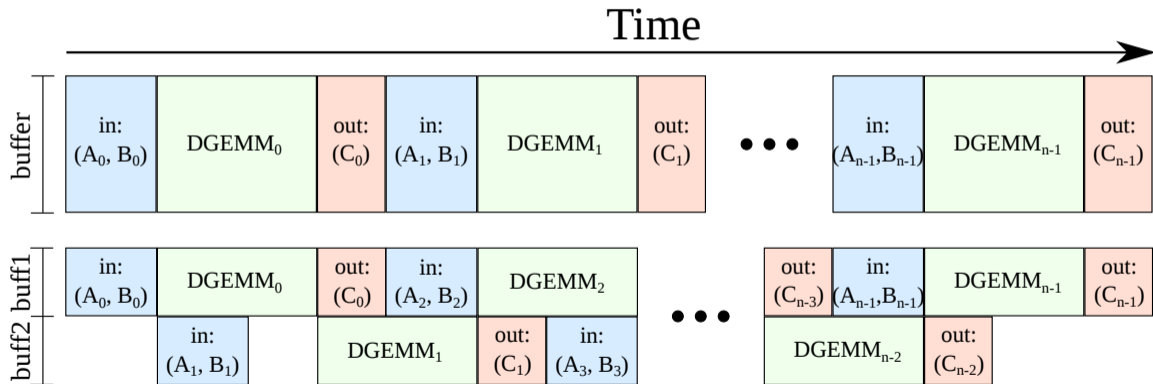
- ▶ By default, `#pragma offload` blocks until offload completes
- ▶ Use clause “signal” with any pointer to avoid blocking
- ▶ Use `#pragma offload_wait` to block where needed

```
1 float* offload0 = &data[0]; // Any unique pointer value as signal
2 #pragma offload target(mic:0) signal(offload0) in(data : length(N))
3 { /* ... will not block code execution because of clause "signal" */ }
4
5 DoSomethingElse();
6
7 /* Now block until offload signalled by pointer "offload0" completes */
8 #pragma offload_wait target(mic:0) wait(offload0)
```

- ▶ Use the target number to avoid hanging

OVERLAPPING COMMUNICATION AND COMPUTATION

Usage example: double buffering to mask communication latency



DOUBLE BUFFERING WITH ASYNCHRONOUS OFFLOAD

```

1  for(int i = 1; i < nMatrices-1; i++) {
2      double* A_T = MatrixA[i]; // Dataset to send next; ...same for B and C
3
4      #pragma offload target(mic:0) signal(A_buff_C) \
5          in(A_buff_C: length(0) alloc_if(0) free_if(0)) ...(same for B and C)
6          { cblas_dgemm(..., A_buff_C, ...); } // Asynchronous offload (COMPUTATION)
7
8      // Send next data set, retrieve previous results (COMMUNICATION):
9      #pragma offload_transfer target(mic:0) in(A_T[0:n*n]: into (A_buff_T[0:n*n]))...
10     #pragma offload_transfer target(mic:0) out(C_buff_T[0:n*n]: into (C_T[0:n*n]))
11     // Wait for asynchronous offload (SYNCHRONIZATION):
12     #pragma offload_wait target(mic:0) wait(A_buff_C)
13
14     if(i%2==1) // Swap Buffers
15         { A_buff_T=A_buff2; A_buff_C =A_buff1; /* ...same for B and C */ }
16     else
17         { A_buff_T=A_buff1; A_buff_C =A_buff2; /* ...same for B and C */ }

```



ADDITIONAL OFFLOAD CONTROLS

TARGET-SPECIFIC CODE

- ▶ During MIC architecture compilation, preprocessor macro `__MIC__` is defined.
- ▶ Allows to fine-tune application performance or output where necessary

```
1 __attribute__((target(mic))) void MyFunction() {  
2 #ifdef __MIC__  
3     printf("I am running on a coprocessor.\n");  
4     const int tuningParameter = 16;  
5 #else  
6     printf("I am running on the host.\n");  
7     const int tuningParameter = 8;  
8 #endif  
9     // ... Proceed, using the variable tuningParameter  
10 }
```

OFFLOAD DIAGNOSTICS

```
vega@lyra% export OFFLOAD_REPORT=2
vega@lyra% ./offload-application
Transferring some data to and from coprocessor...
Done. Bye!
[Offload] [MIC 0] [File]                offload-application.cc
[Offload] [MIC 0] [Line]                6
[Offload] [MIC 0] [CPU Time]            0.505982 (seconds)
[Offload] [MIC 0] [CPU->MIC Data]      1024 (bytes)
[Offload] [MIC 0] [MIC Time]           0.000409 (seconds)
[Offload] [MIC 0] [MIC->CPU Data]      1024 (bytes)
vega@lyra%
```

- ▶ Set environment variable `OFFLOAD_REPORT` to 1 or 2 for automatic collection and output of offload information.
- ▶ Unset or set `OFFLOAD_REPORT=0` to disable offload diagnostics

OFFLOAD DEVICES, SPECIFYING AVAILABLE COPROCESSORS

- ▶ Specify coprocessors to use; For example (using 0 and 1),

```
vega@lyra% export OFFLOAD_DEVICES=0,1
```

- ▶ Disable Offloading

```
vega@lyra% export OFFLOAD_DEVICES=none
```

Disabling Offload is useful for debugging. For example;

```
vega@lyra% icpc Offload-Fallback.cc -o Offload-Fallback
vega@lyra% ./Offload-Fallback
Hello from offload on MIC with 244 logical processors.
vega@lyra% export OFFLOAD_DEVICES=none # Coprocessors disabled
vega@lyra% ./Offload-Fallback
Hello from offload on CPU with 48 logical processors.
```

ENVIRONMENT VARIABLE FORWARDING WITH OFFLOAD

- ▶ By default, all host environment variables on the host will be copied to the coprocessor when offload starts.
- ▶ In order to have different values for an environment variable on host and coprocessor, set `MIC_ENV_PREFIX`
- ▶ The prefix is dropped when variables are copied to coprocessor

```
vega@lyra% # This sets the value of OMP_NUM_THREADS on the host:
vega@lyra% export OMP_NUM_THREADS=48
vega@lyra%
vega@lyra% # This enables special rules for variable copying:
vega@lyra% export MIC_ENV_PREFIX=XEONPHI
vega@lyra%
vega@lyra% # This sets the value of OMP_NUM_THREADS on the coprocessor:
vega@lyra% export XEONPHI_OMP_NUM_THREADS=240
```



OFFLOAD IN OPENMP 4.0

OPENMP 4.0 TARGET OFFLOAD

- ▶ Another API for offload: `#pragma omp target`
- ▶ Part of the OpenMP 4.0 standard
- ▶ Designed as portable solution (coprocessors, GPGPUs)
- ▶ On Xeon Phi, uses the same back-end as `#pragma offload`

```
1 #pragma omp target
2 {
3 #pragma omp parallel for
4   for(int i=0; i<size; i++)
5     data[i] = 0;
6 }
```

Application runs on the host, but some parts of code and data are moved (“offloaded”) the coprocessor. Scope-local scalars and stack arrays offloaded automatically.

CLAUSES OF PRAGMA OMP TARGET

```
1 #pragma omp target [clause[, clause[, ...]]]
```

- ▶ `device(int)` – offload to a specific device (coprocessor)
- ▶ `map([type:] variables)` – create data environment. `type` is `to`, `from`, `tofrom` or `alloc`
- ▶ `if(expr)` – optional offload

Link to [reference manual](#).

OPENMP 4.0 TARGET DATA MAPPING

Use `#pragma omp target data` to create a device data environment. This allows to keep persistent data on coprocessor. Example:

```
1 #pragma omp target data map(from:data)
2 {
3 #pragma omp target
4 #pragma omp parallel for
5     for(int i=0; i<size; i++) data[i] = 0;
6
7 #pragma omp target
8 #pragma omp parallel for
9     for(int i=0; i<size; i++) data[i] += 1;
10 }
```

data array copied back from coprocessor only once at the end.
Link to [reference manual](#).

MOVEMENT OF PERSISTENT DATA

Use `#pragma omp target update` to force data movement within the data environment. Example:

```
1 #pragma omp target data map(from:data)  
2 {  
3 #pragma omp target  
4   { ... }  
5  
6 #pragma omp target update from(data)  
7  
8 #pragma omp target  
9   { ... }  
10 }
```

data array copied from coprocessor between offloads, and at the end.

Link to [reference manual](#).

OFFLOADING FUNCTIONS WITH `#pragma omp target`

Use `#pragma omp declare target` on functions that may be offloaded (similar to `__attribute__((target(mic)))`). Example:

```
1  #pragma omp declare target
2  void myinit(int* data, int size){
3  #pragma omp parallel for
4      for(int i=0; i<size; i++) data[i] = 0;
5  }
6  #pragma omp end declare target
7
8  int main(int argv, char** argc){
9      ...
10 #pragma omp target map(tofrom:data) map(to:size)
11     myinit(data, size);
12 }
```

Link to [reference manual](#).

#pragma offload target **VS.** #pragma omp target

1. Different interfaces to the same offload library back-end
2. #pragma offload target is Intel-specific, #pragma omp target is part of a cross-platform standard (although, as today, cross-platform support is not widespread).
3. #pragma offload target allows data/memory persistence outside of the scope of a pragma, #pragma omp target – only within the lexically structured scope (correct us if we're wrong).
4. #pragma offload is a more flexible model and will continue to be supported (see Intel's [communication](#)).

Additional information: [webinar](#).



SHARED VIRTUAL MEMORY OFFLOAD MODEL

SHARED VIRTUAL MEMORY MODEL

```
1  _Cilk_shared int arr[N]; // This is a virtual-shared array
2
3  _Cilk_shared void Compute() { // This function may be offloaded
4      // ... function uses array arr[]
5  }
6
7  int main() {
8      // arr[] can be initialized on the host
9      _Cilk_offload Compute(); // and used on coprocessor
10     // and the values are returned to the host
11 }
```

- ▶ Alternative to Explicit Offload
- ▶ Data synced from host to coprocessor before the start of offload
- ▶ Data synced from coprocessor to host at the end of offload

SHARED VIRTUAL MEMORY MODEL

```
1  int* _Cilk_shared data; // Pointer to a virtual-shared array
2
3  int main() {
4      // Working with pointer-based data is illustrated below:
5      data = (_Cilk_shared int*)_Offload_shared_malloc(N*sizeof(float));
6      _Offload_shared_free(data);
7  }
```

- ▶ Addresses of virtual-shared pointers identical on host and coprocessors
- ▶ Synchronized before and after offload, with page granularity

REVIEW AND WHAT'S NEXT

- ▶ Native programming = treat coprocessor as a processor
- ▶ Offload = launched on host, some functions on coprocessor
- ▶ Explicit offload – `#pragma offload` or `#pragma omp target` – allows data marshalling; for bitwise-copyable data
- ▶ Shared virtual memory – `_Cilk_shared/_Cilk_offload` – automatic coherence; for complex objects

Next session: expressing data parallelism, vectorization.



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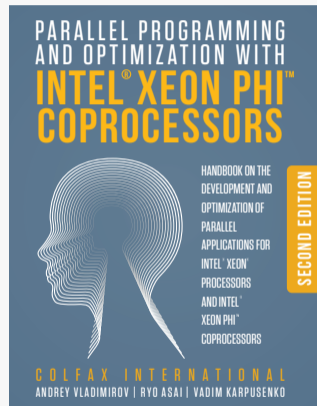
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Episode 2.1 — Purpose of the MIC architecture



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